

Polymorphic Symmetric Multiple Dispatch with Variance

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1 Background

1.1 Method Overloading

- Allows defining multiple methods with the same name.

```
add(x: Int, y: Int): Int = ...
add(x: List, y: List): List = ...
add(x: SortedList, y: List): SortedList = ...
add(x: SortedList, y: SortedList): SortedList = ...
```

- A method dispatch mechanism chooses a method to invoke.

```
a: List = SortedList(1, 3, 5)
b: List = SortedList(2, 4, 6)           (SortedList <: List)
add(a, b)
```

1.2 Static Dispatch

- Selects a method based on the compile-time types of arguments.

```
add(x: List, y: List): List = ...
```

1.3 Single Dispatch

- Considers the run-time type of a single argument, usually a receiver.

```
add(x: SortedList, y: List): SortedList = ...
```

- Is widely used in modern object-oriented languages.
- Causes the binary method problem.

1.4 Multiple Dispatch

- Considers the run-time types of more than one argument.
- Can be either symmetric or asymmetric.

```
add(x: SortedList, y: SortedList): SortedList = ...
```

- Provides an intuitive solution for the binary method problem.

2 Motivation

2.1 The Goals of Method Dispatch

- Every method invocation should be unambiguous.
- The most specific declaration exists among applicable declarations.

```
add(x: List, y: List): List = ...
add(x: SortedList, y: List): SortedList = ...
add(x: SortedList, y: SortedList): SortedList = ...
```

more specific

- Every method invocation should be type-sound.
- A dynamically chosen method preserves the return type of a statically chosen method.

```
add(x: List, y: List): List = ...
add(x: SortedList, y: SortedList): SortedList = ...
```

2.2 Issues in Symmetric Multiple Dispatch

- Run-time types are unknown at compile time.
- Cannot check unambiguity and type preservation statically.

```
add(x: List, y: SortedList): SortedList = ...
add(x: SortedList, y: List): SortedList = ...
```

```
a: SortedList = SortedList(1, 3, 5)
b: List = SortedList(2, 4, 6)
add(a, b)
```

- At compile time, it is unambiguous, but not at run time.
- Needs a disambiguating declaration.

```
add(x: SortedList, y: SortedList): List = ...
```

- It does not preserve static return type.
- Needs a fix for the return type.

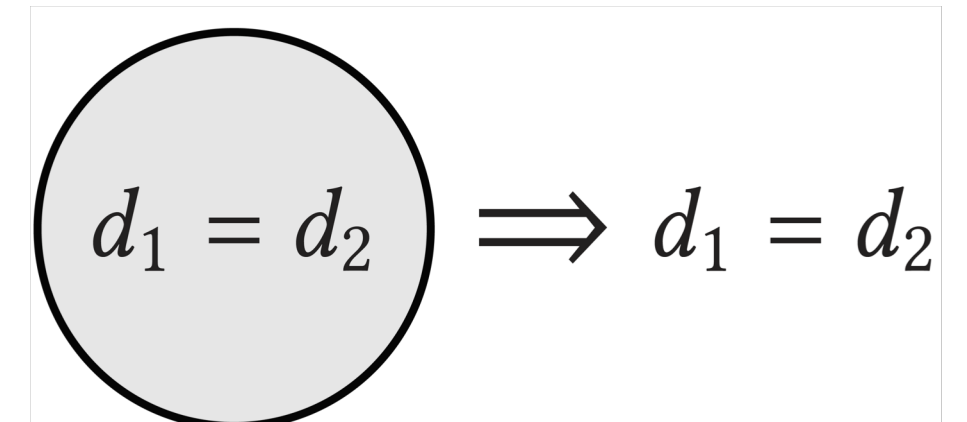
```
add(x: SortedList, y: SortedList): SortedList = ...
```

- We can check validities of overloaded method sets statically.

3 Overloading Rules

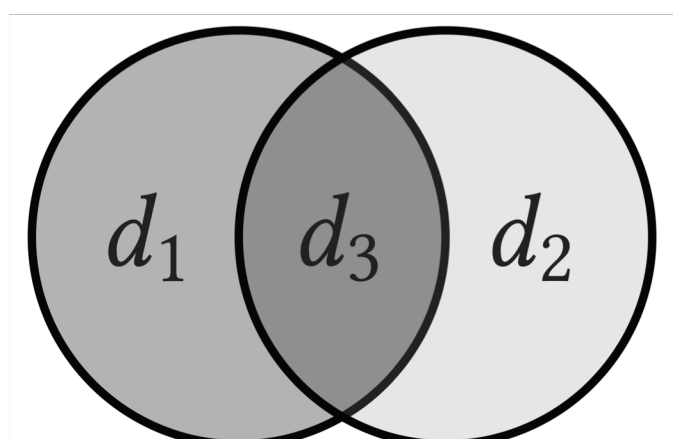
3.1 No Duplicate Rule

- Rules out duplicate signatures.



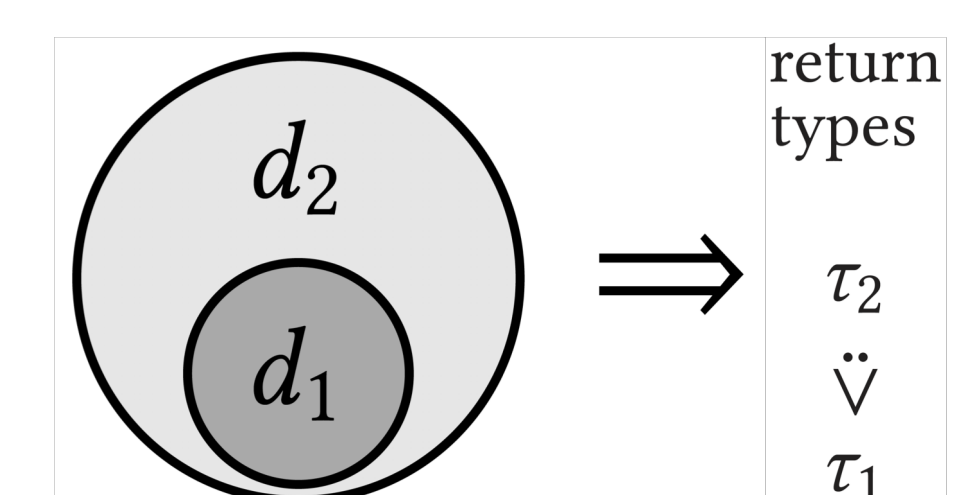
3.2 Meet Rule

- Ensures existence of the most specific declaration.



3.3 Return Type Rule

- Ensures type preservation.



4 Additional Issues in the Paper

4.1 Parametric Polymorphism

- Parametrizes classes and methods by types.

```
add[P](x: List[P], y: List[P]): List[P] = ...
```

- We introduce quantified types to formally define "more specific".
- Variance defines additional subtype relation for parametric types.
 - If T is covariant, then $A <: B \rightarrow T[A] <: T[B]$.
 - If T is contravariant, then $A <: B \rightarrow T[B] <: T[A]$.

4.2 Dynamic Dispatch Algorithm

- Determines a method to invoke.
- Infers type arguments for parametric methods.
- We formally defined the algorithm and proved its correctness.